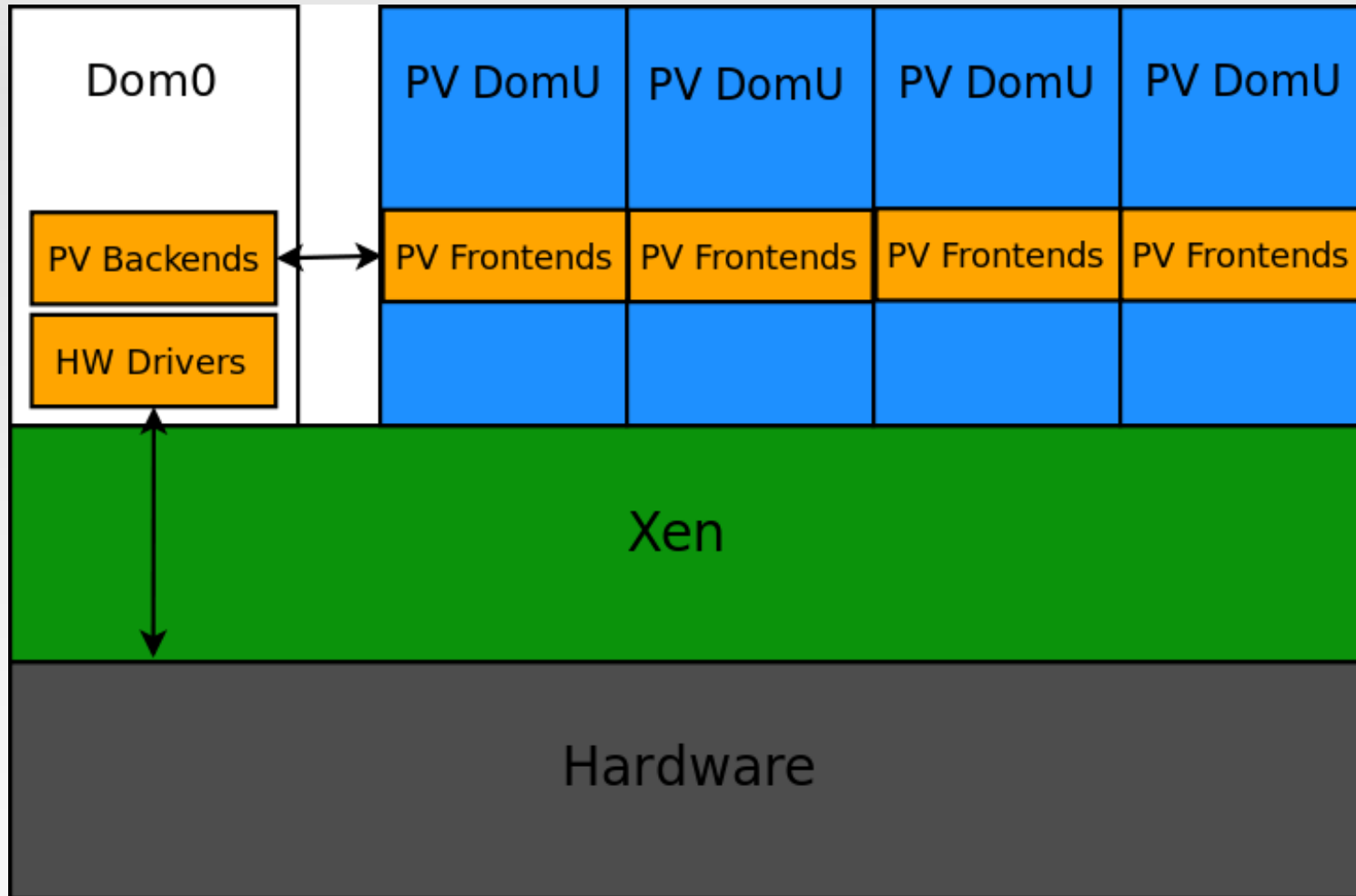


# Xen

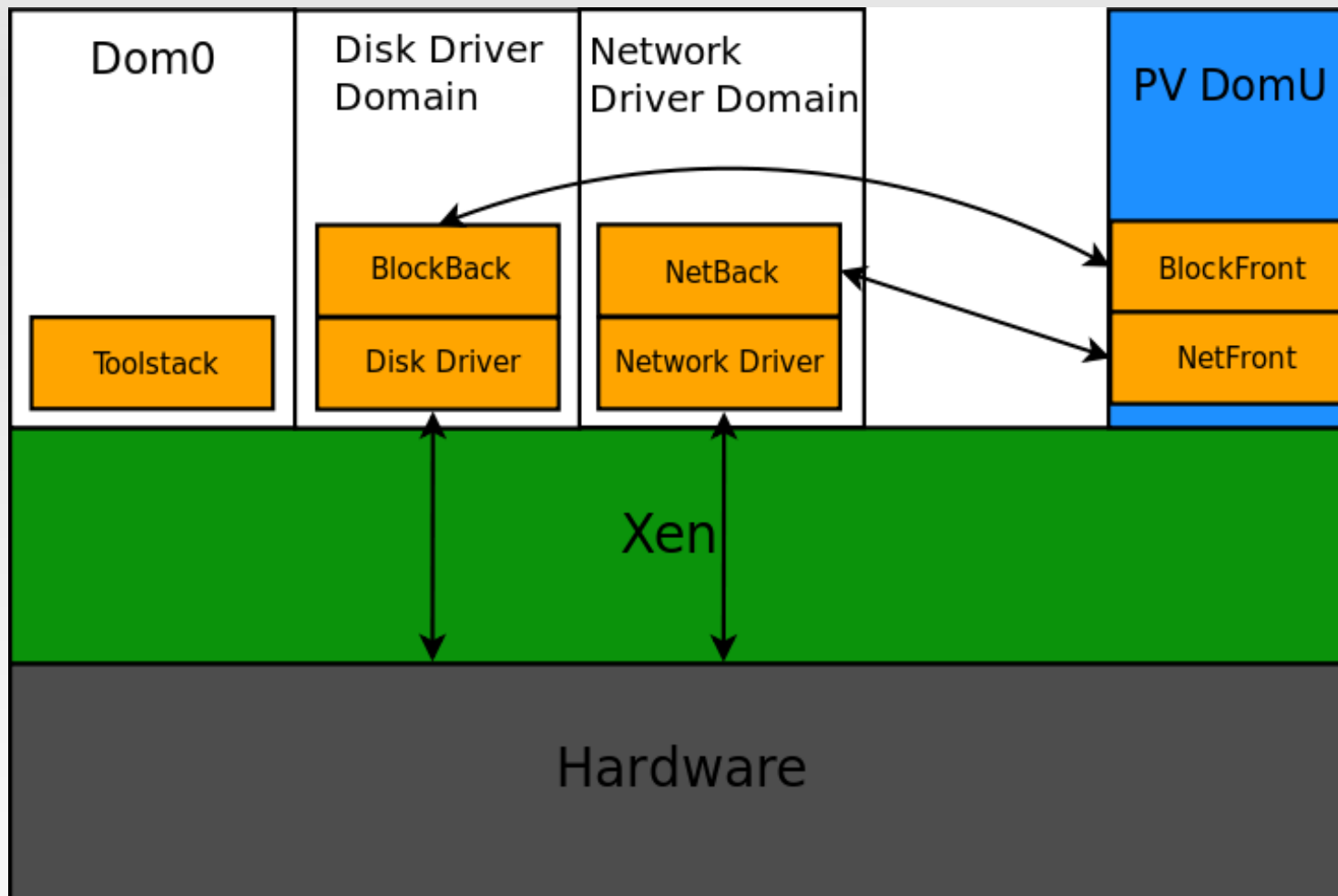
past, present and future

**Stefano Stabellini**

# Xen architecture: PV domains



# Xen arch: driver domains



# Xen: advantages

- small surface of attack
- isolation
- resilience
- specialized algorithms (scheduler)

# Xen and the Linux kernel

Xen was initially a university research project



invasive changes to the kernel to run Linux as a  
PV guest

even more changes to run Linux as dom0

# Xen and the Linux kernel

Xen support in the Linux kernel not upstream



Great maintenance effort on distributions



Risk of distributions dropping Xen support

# Xen and the Linux kernel

- PV support went in Linux 2.6.26
- basic Dom0 support went in Linux 2.6.37
- Netback went in Linux 2.6.39
- Blkback went in Linux 3.0.0

A single 3.0.0 Linux kernel image boots on native, on Xen as domU, as dom0 and PV on HVM guest

# Xen and Linux distributions

## *2010*

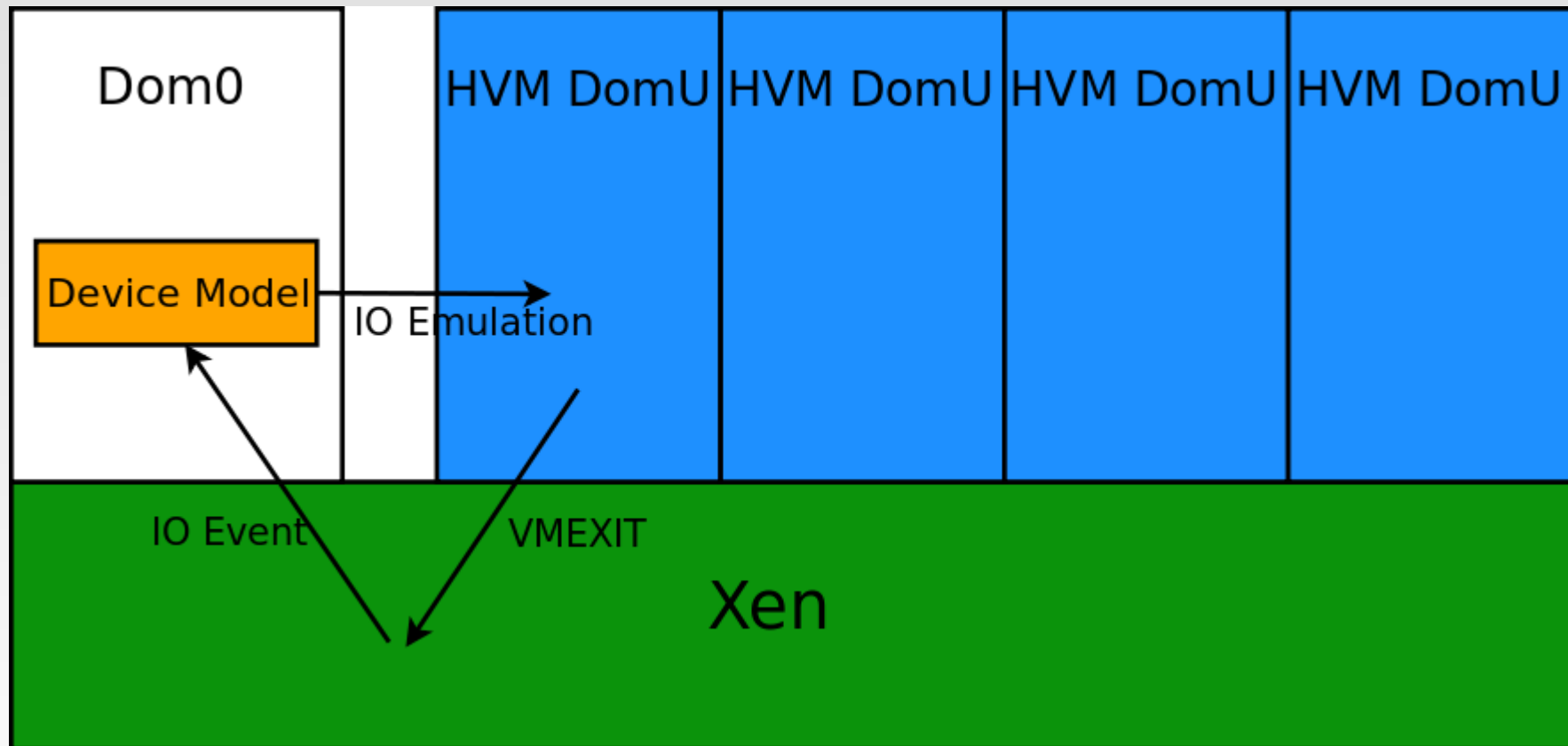
- Fedora and Ubuntu dropped Xen support from their Linux kernels
- Debian, Suse, Gentoo still provide Xen kernels
- XenServer went Open Source with XCP

## *Present*

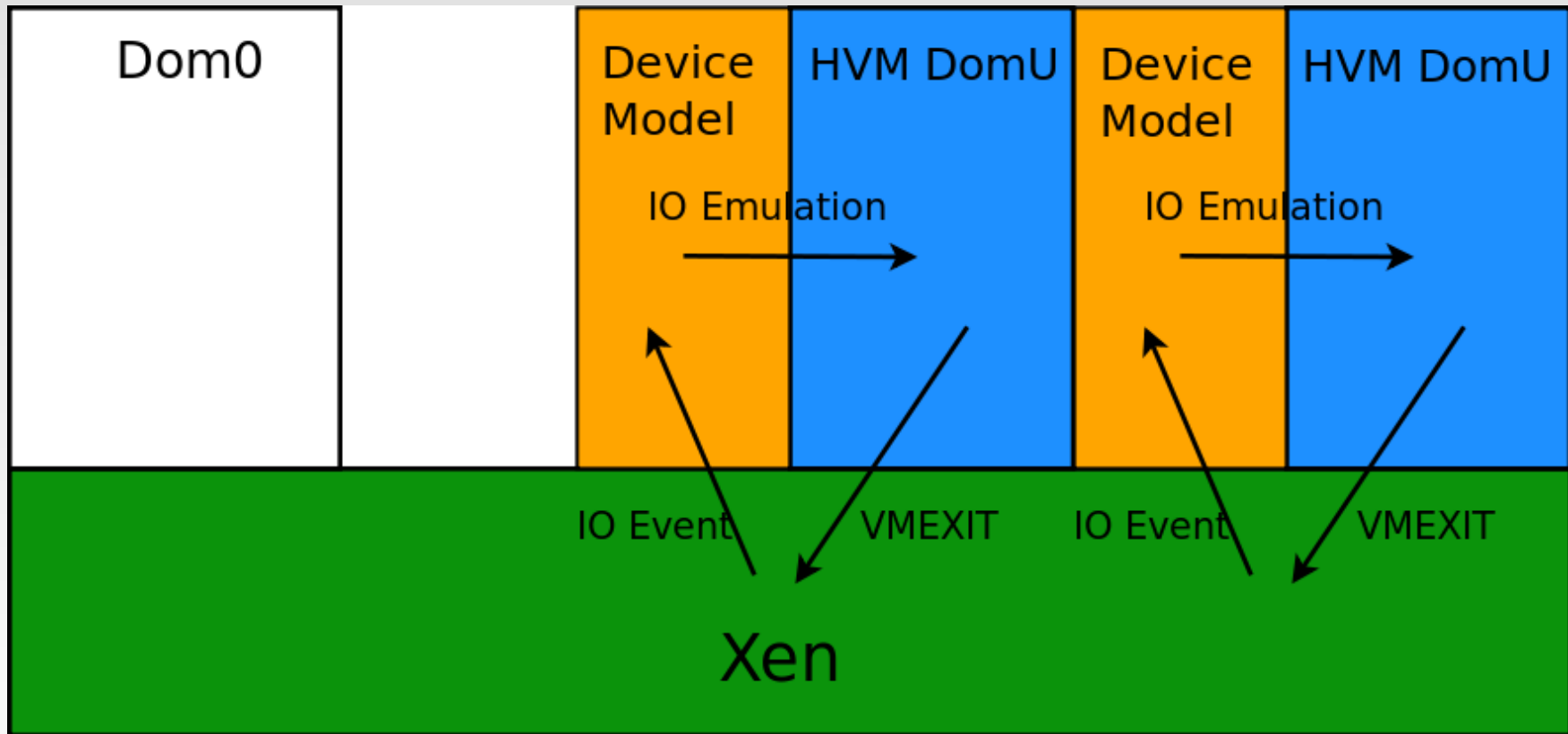
- Fedora and Ubuntu are adding Xen support back in kernel in the next releases



# Xen architecture: HVM domains



# Xen architecture: stubdoms



# Xen and Qemu

- initially forked in 2005
- updated once every few releases
- Xen support went in upstream Qemu at the beginning of 2011
- Upstream Qemu is going to be used as device model with Xen 4.2

# New developments: Libxenlight

Multiple toolstacks:

- Xend, Xapi, XenVM, LibVirt, ...
- code duplications, inefficiencies, bugs, wasted efforts

Xend:

- difficult to understand, modify and extend
- significant memory footprint

# Libxenlight

What is Libxenlight:

- a small lower level library in C
- simple to understand
- easy to modify and extend

Goals:

- provide a simple and robust API for toolstacks
- create a common codebase to do Xen operations

# XL

- the unit testing tool for libxenlight
- feature complete
- a minimal toolstack
- compatible with xm

Do more with less!

# XL: design principles

- smallest possible toolstack on top of libxenlight
- stateless

CLI → XL → libxenlight → EXIT

# XL vs. Xend

## XL: pros

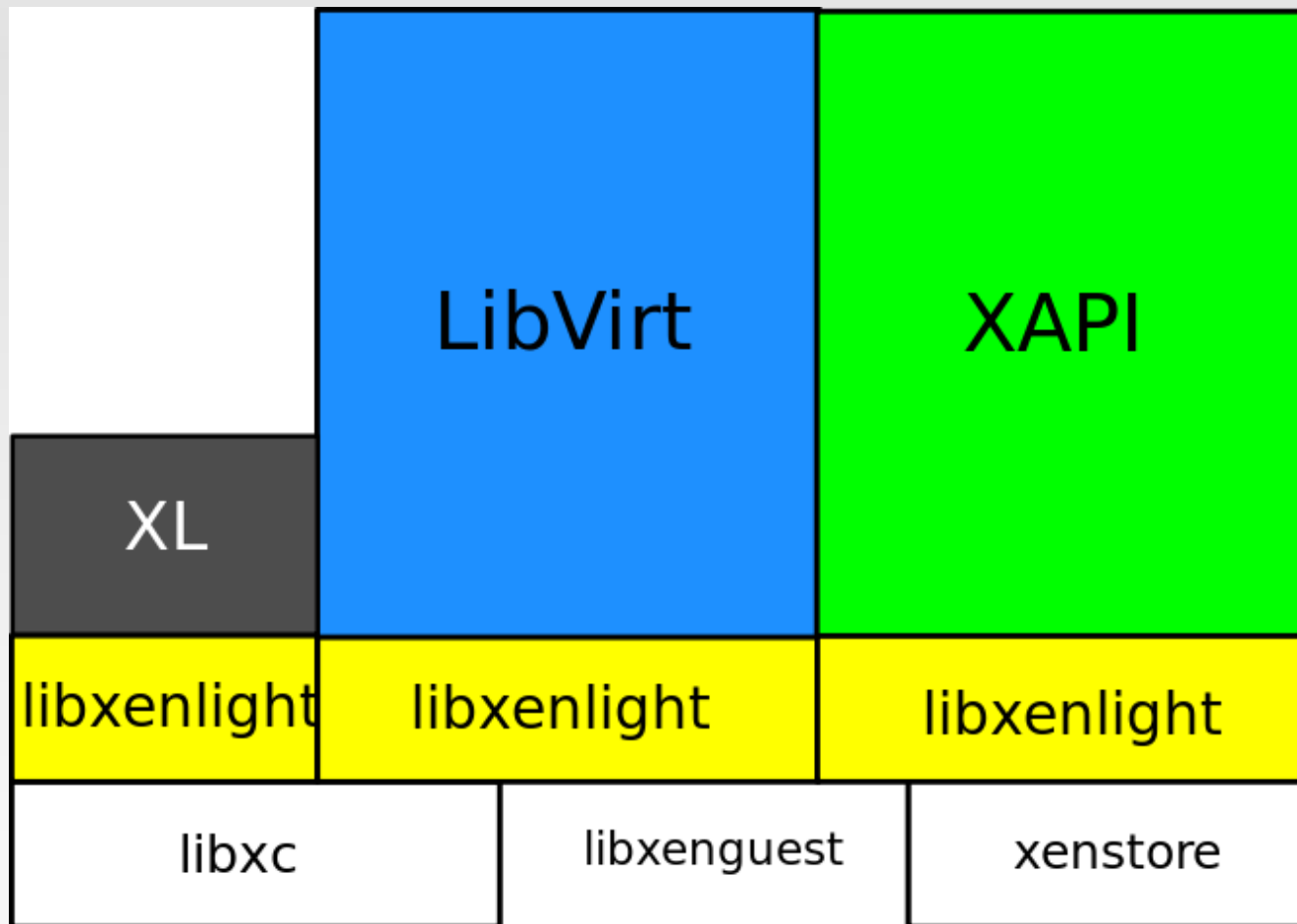
- very small and easy to read
- well tested
- compatible with xm

## Xend: pros

- provide XML RPC interface
- provide "managed domains"



# Libxenlight: the new world



# Linux PV on HVM

paravirtualized interfaces in HVM guests

# Linux as a guests: problems

Linux PV guests have limitations:

- difficult “different” to install
- limited set of virtual hardware

Linux HVM guests:

- install the same way as native
- very slow

# Linux PV on HVM: the solution

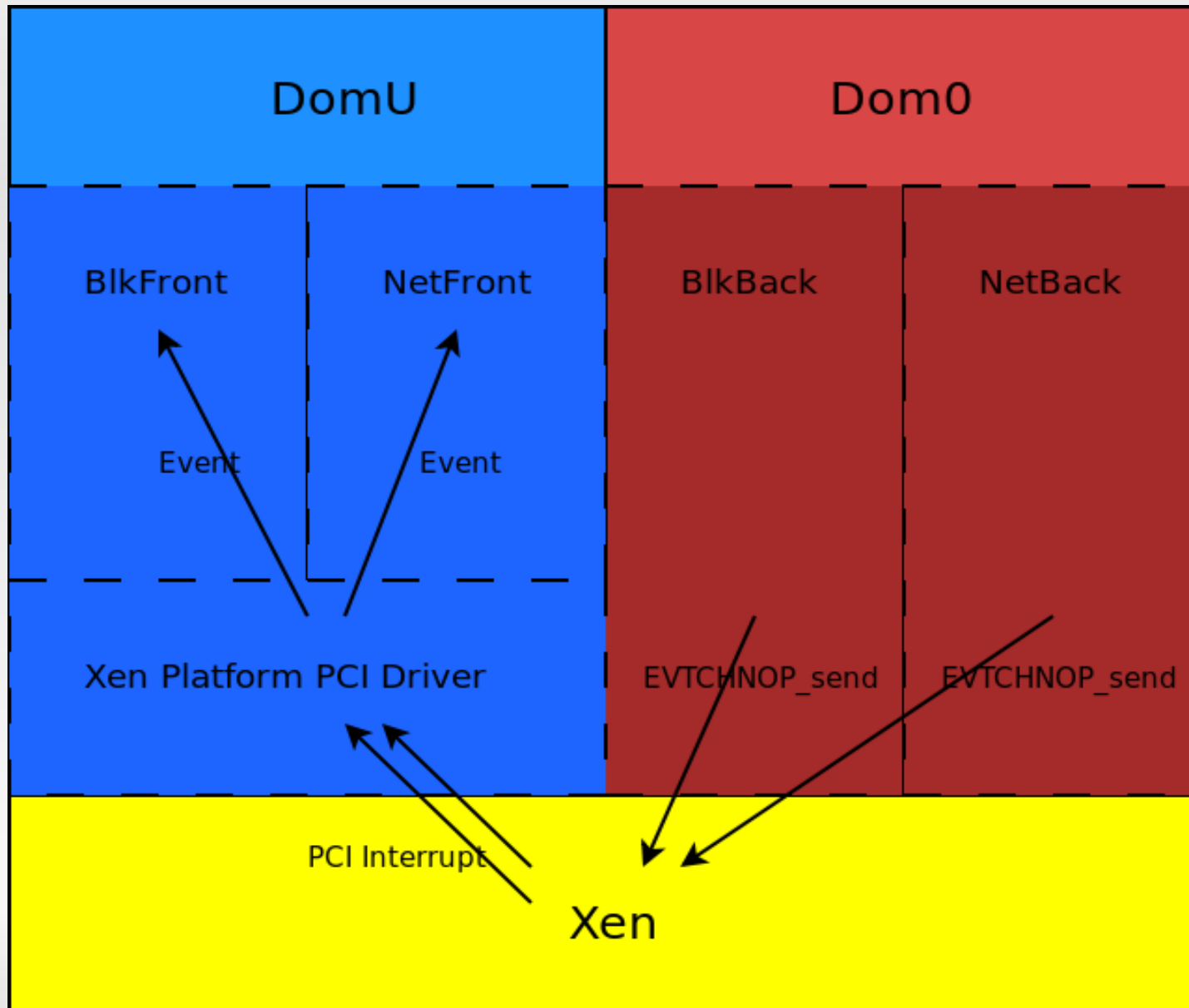
- install the same way as native
- PC-like hardware
- access to fast paravirtualized devices
- exploit nested paging

# Linux PV on HVM: initial feats

Initial version in Linux 2.6.36:

- introduce the xen platform device driver
- add support for HVM hypercalls, xenbus and grant table
- enables **blkfront**, **netfront** and **PV timers**
- add support to PV suspend/resume
- the **vector callback** mechanism

# Old style event injection



# Receiving an interrupt

do\_IRQ

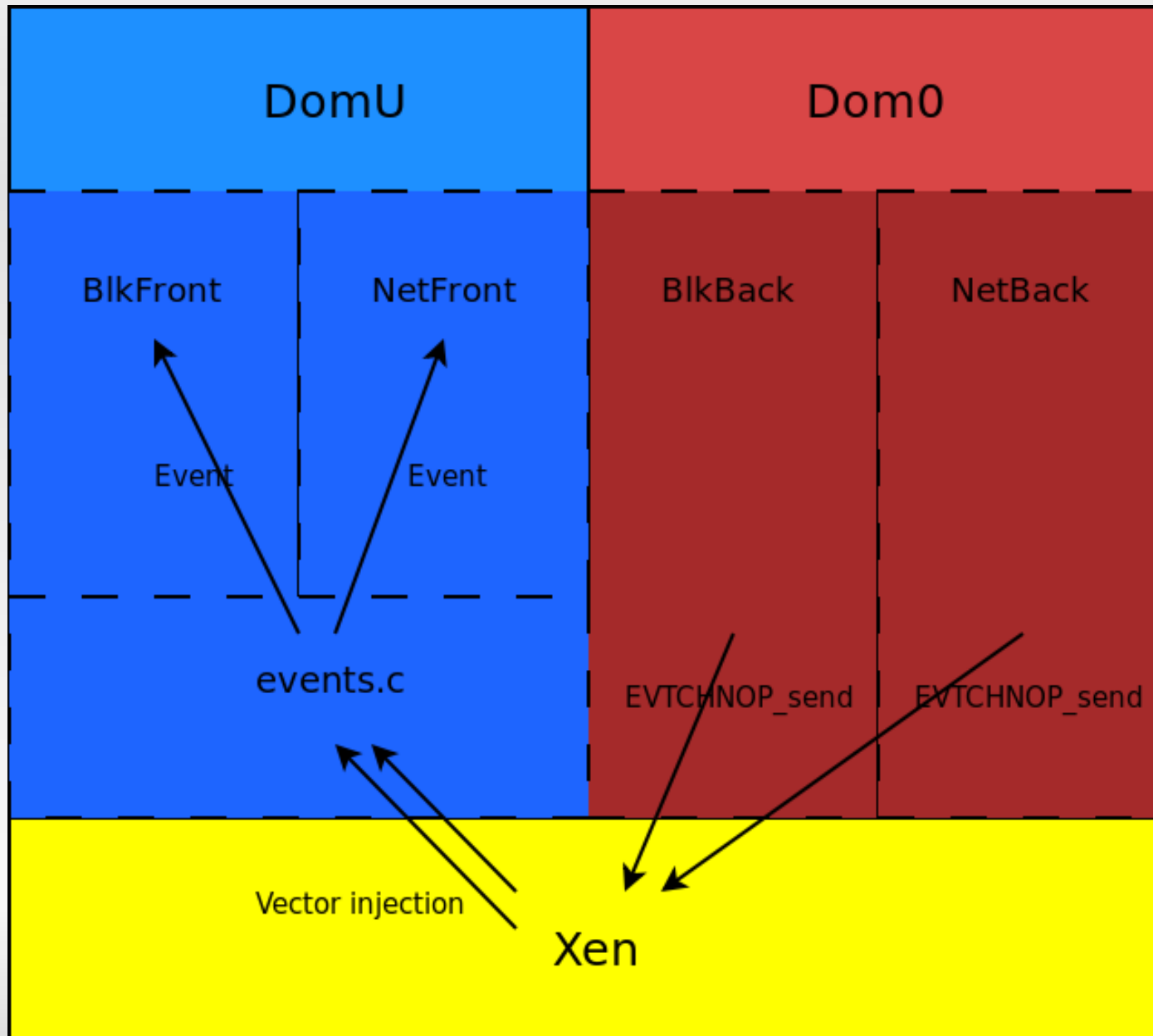
    handle\_fasteoi\_irq

        handle\_irq\_event

            xen\_evtchn\_do\_upcall

**ack\_apic\_level ← ≥3 VMEXIT**

# The new vector callback





# Receiving a vector callback

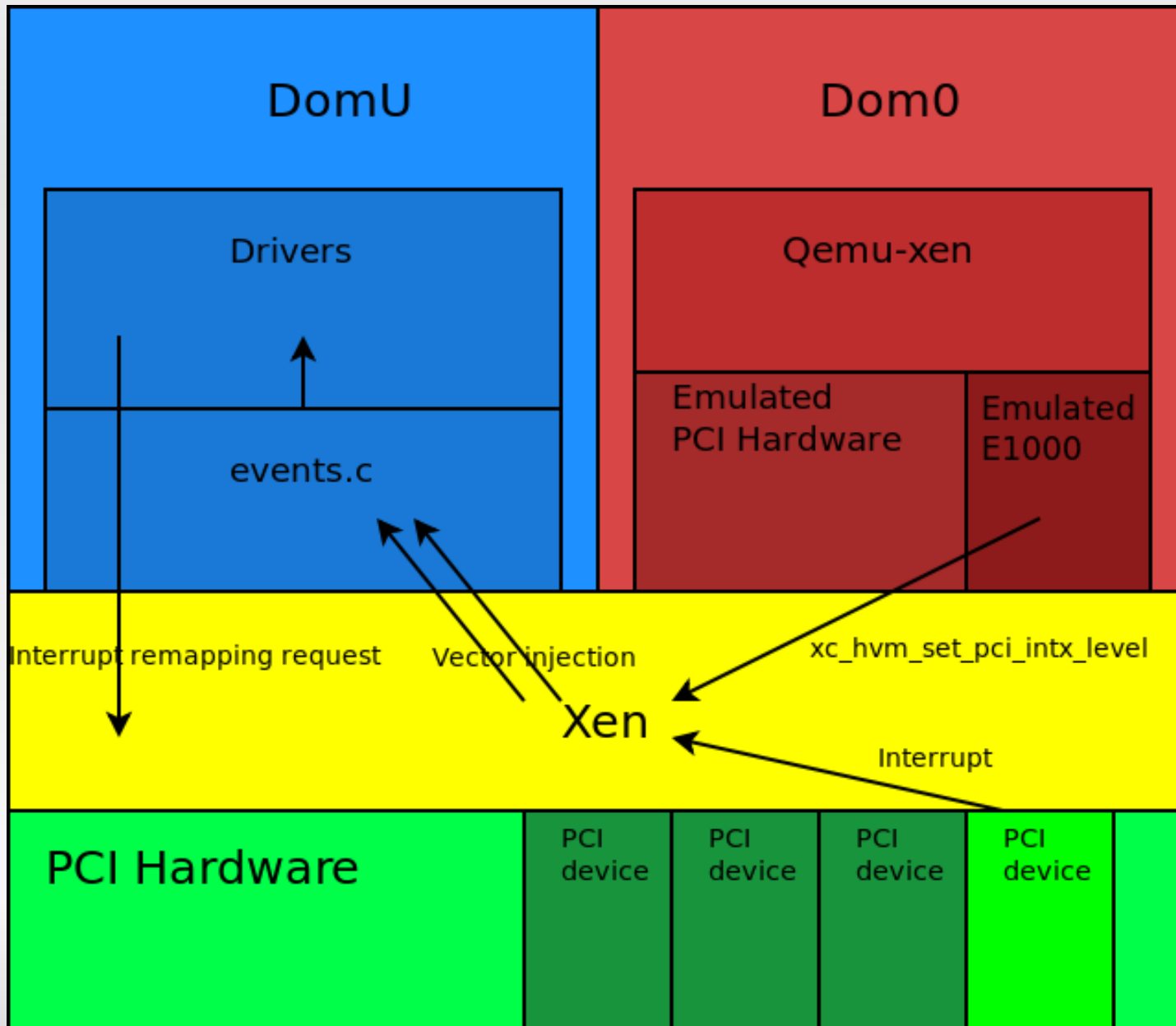
xen\_evtchn\_do\_upcall

# Linux PV on HVM: newer feats

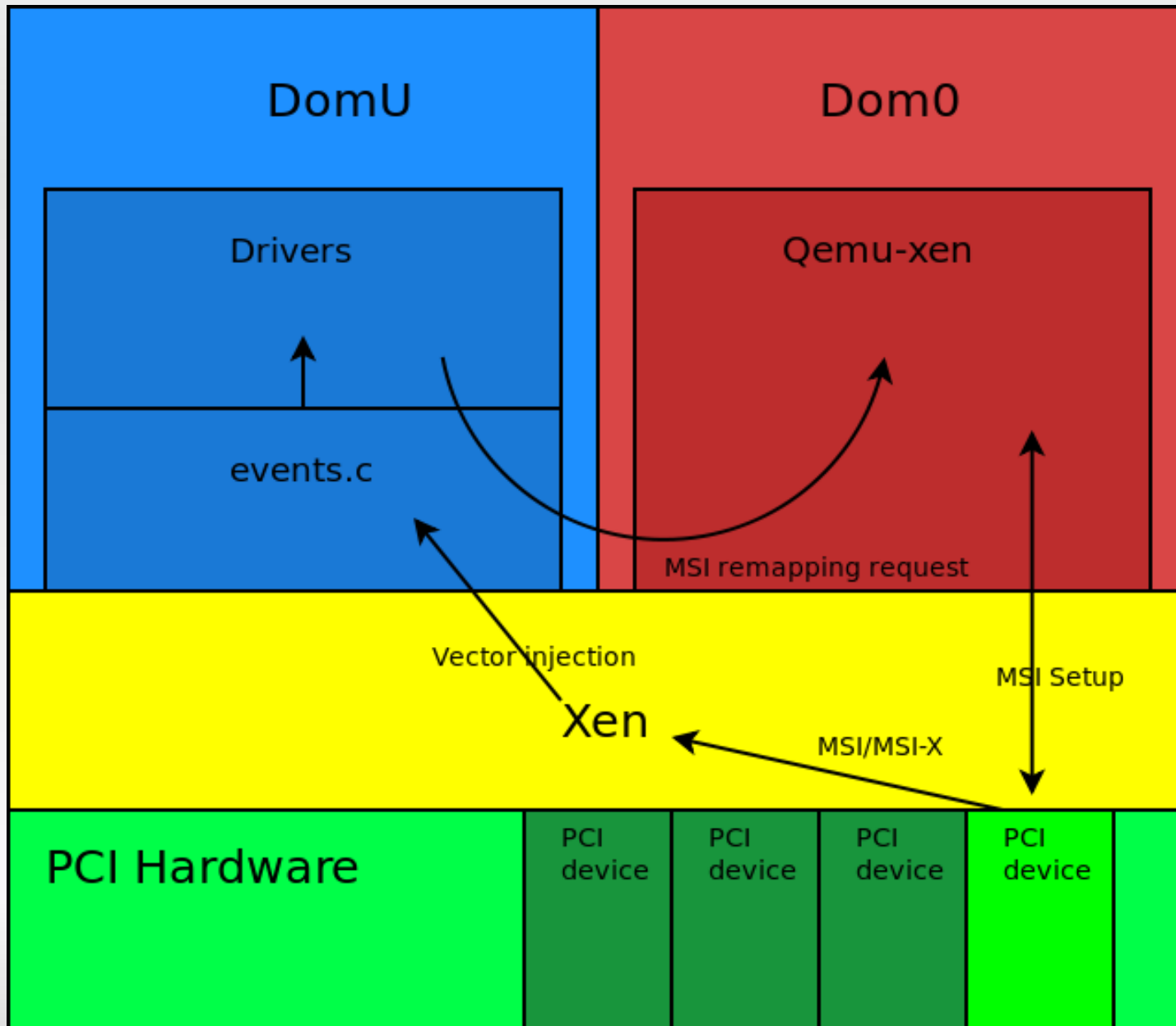
Later enhancements (2.6.37+):

- ballooning
- PV spinlocks
- PV IPIs
- Interrupt remapping onto event channels
- MSI remapping onto event channels

# Interrupt remapping



# MSI remapping



# PV spectrum

	HVM guests	Classic PV on HVM	Enhanced PV on HVM	Hybrid PV on HVM	PV guests
Boot sequence	emulated	emulated	emulated		paravirtualized
Memory	hardware	hardware	hardware		paravirtualized
Interrupts	emulated	emulated	paravirtualized		paravirtualized
Timers	emulated	emulated	paravirtualized		paravirtualized
Spinlocks	emulated	emulated	paravirtualized		paravirtualized
Disk	emulated	paravirtualized	paravirtualized		paravirtualized
Network	emulated	paravirtualized	paravirtualized		paravirtualized
Privileged operations	hardware	hardware	hardware		paravirtualized

# Benchmarks: the setup

## Hardware setup:

Dell PowerEdge R710

CPU: dual Intel Xeon E5520 quad core CPUs @ 2.27GHz

RAM: 22GB

## Software setup:

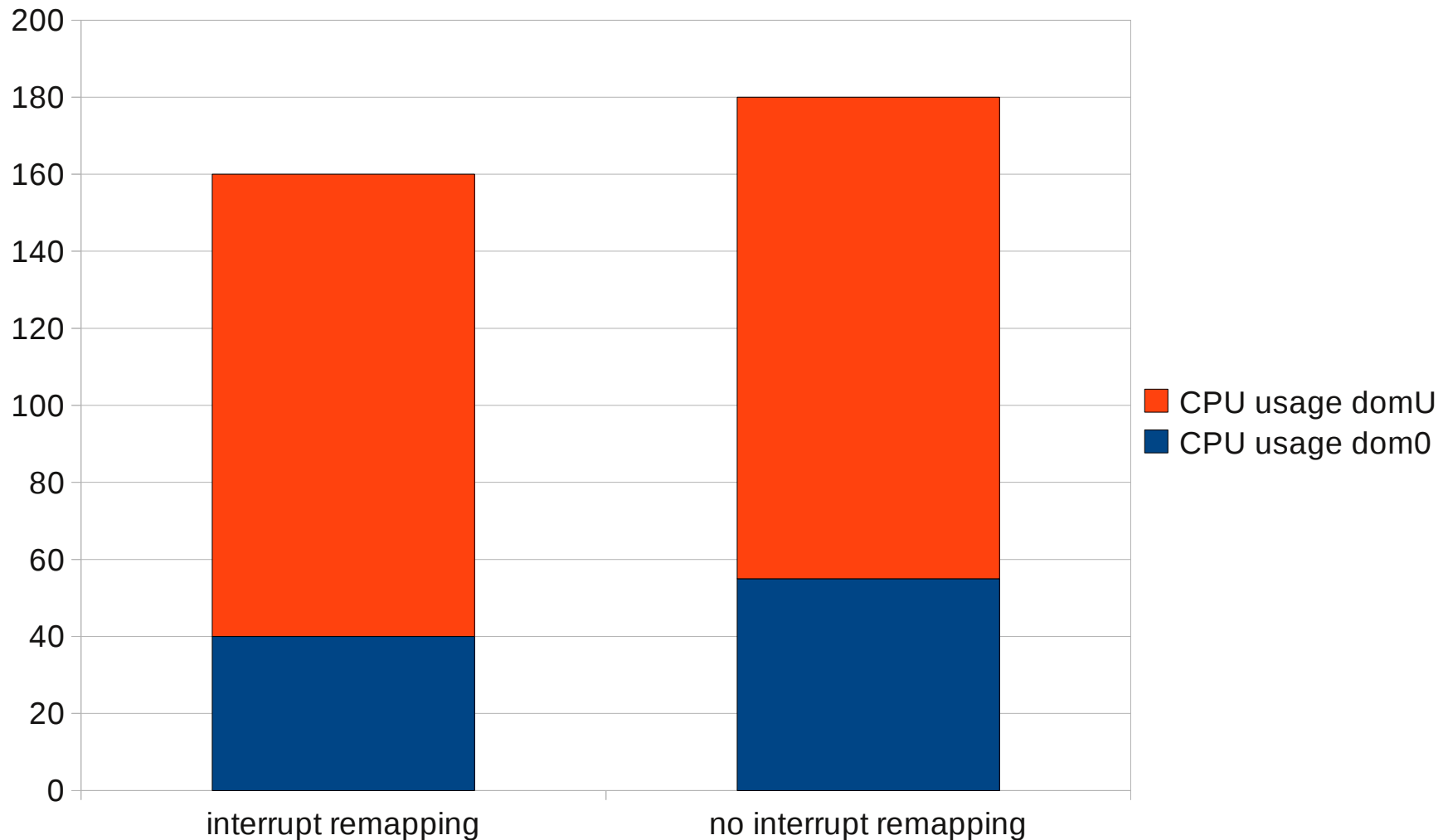
Xen 4.1, 64 bit

Dom0 Linux 2.6.32, 64 bit

DomU Linux 3.0 rc4, 8GB of memory, 8 vcpus

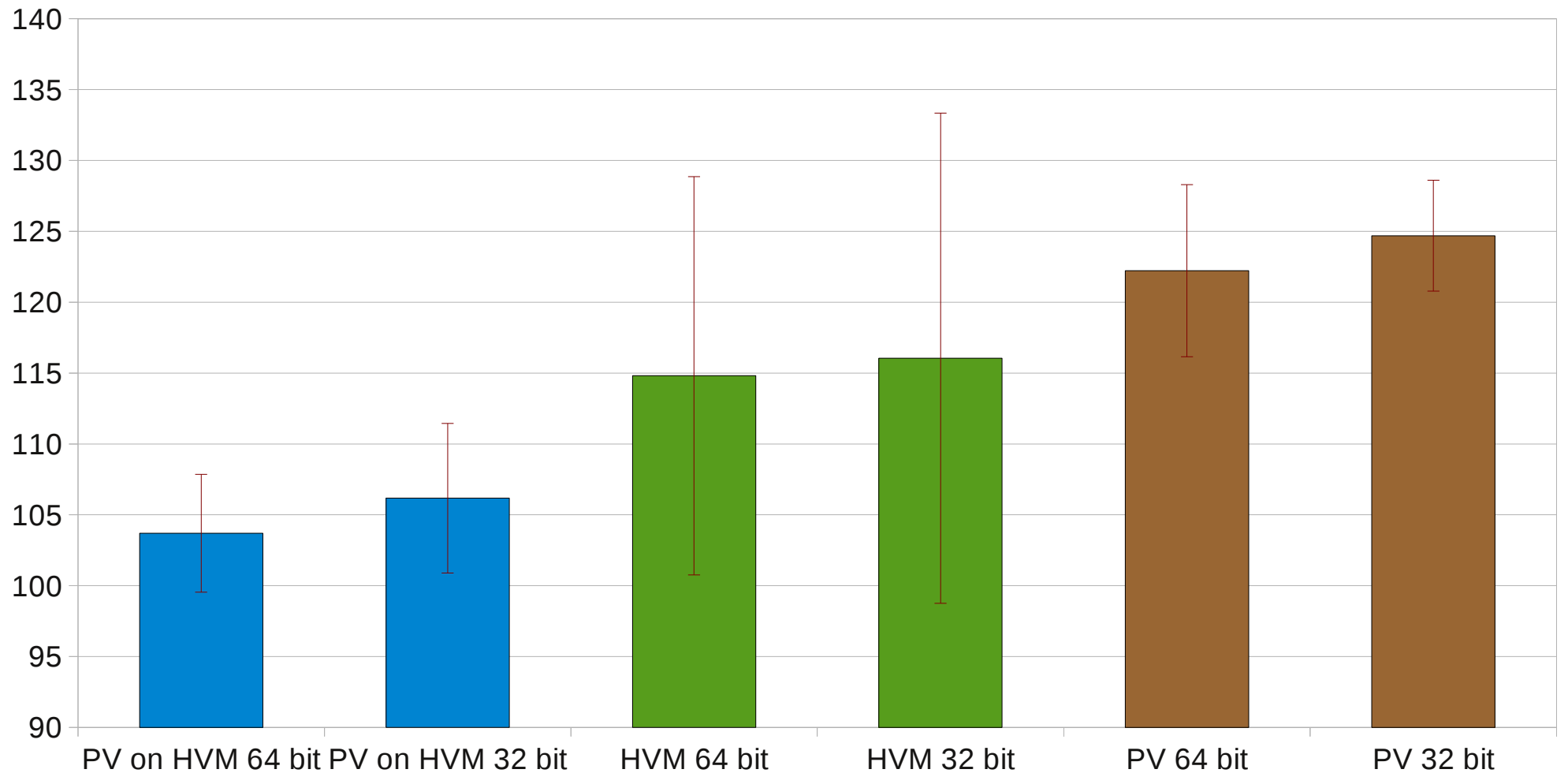
# PCI passthrough: benchmark

PCI passthrough of an Intel Gigabit NIC  
CPU usage: the lower the better:



# Kernbench

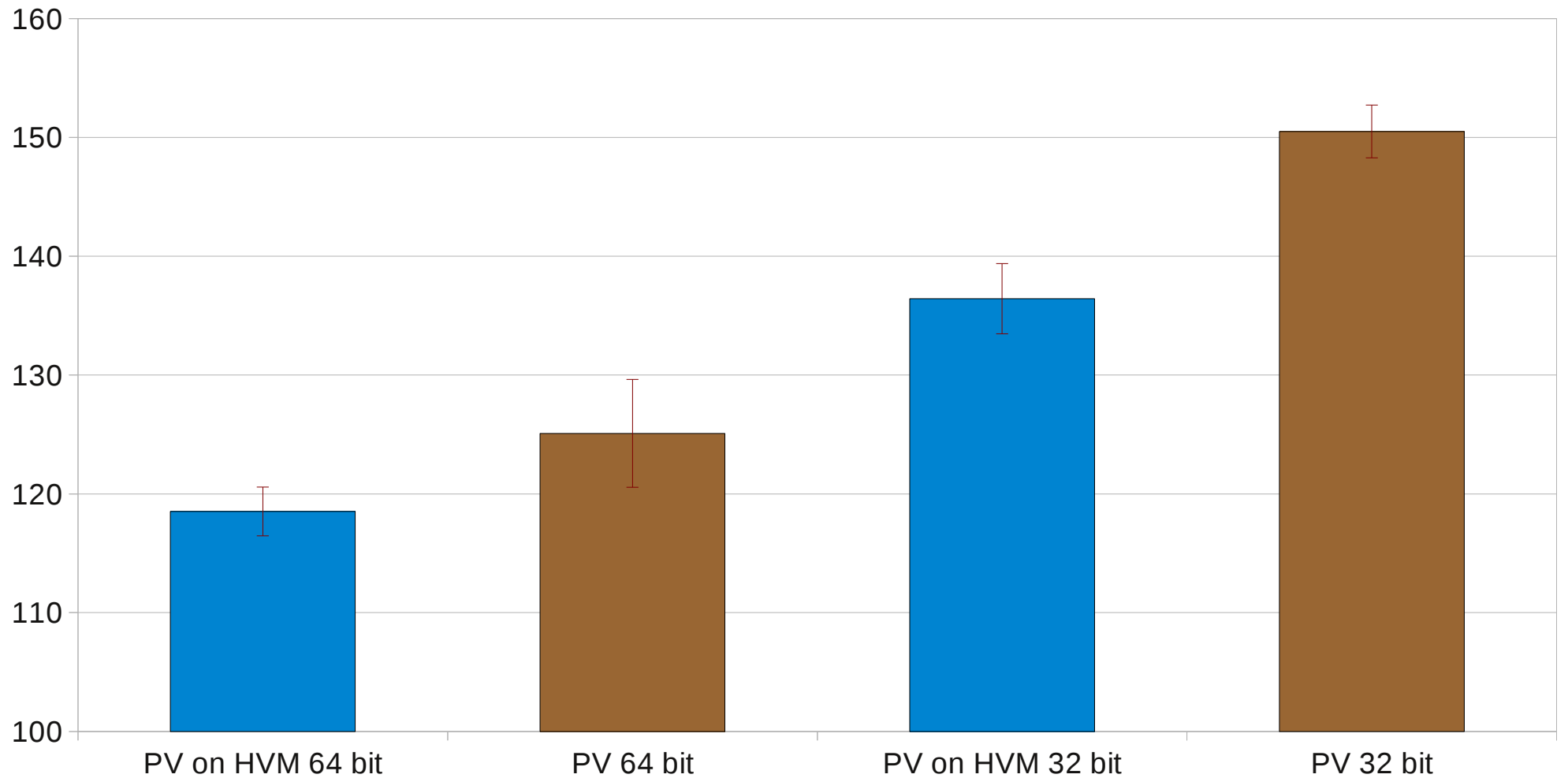
Results: percentage of native, the lower the better





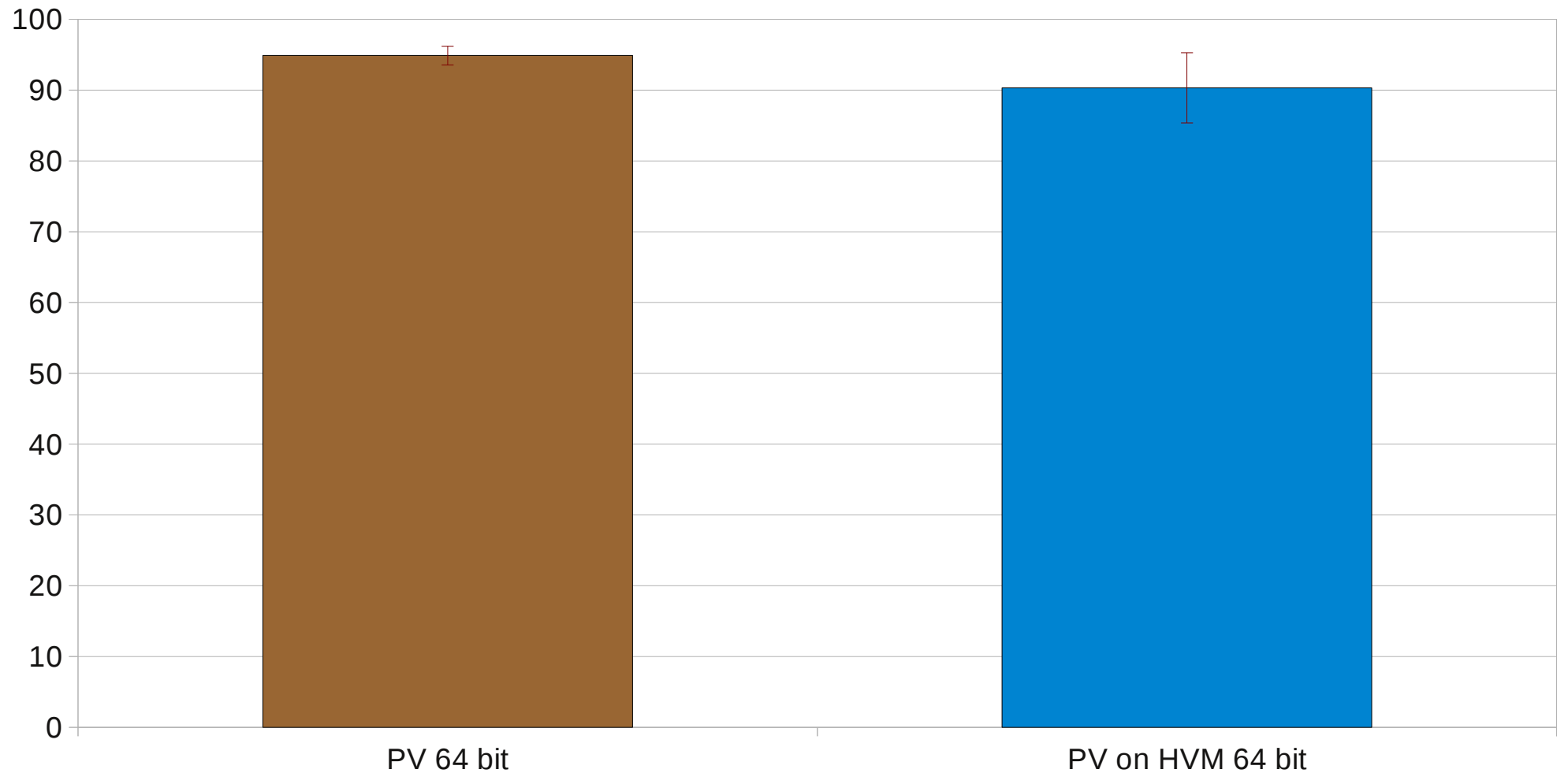
# PBZIP2

Results: percentage of native, the lower the better



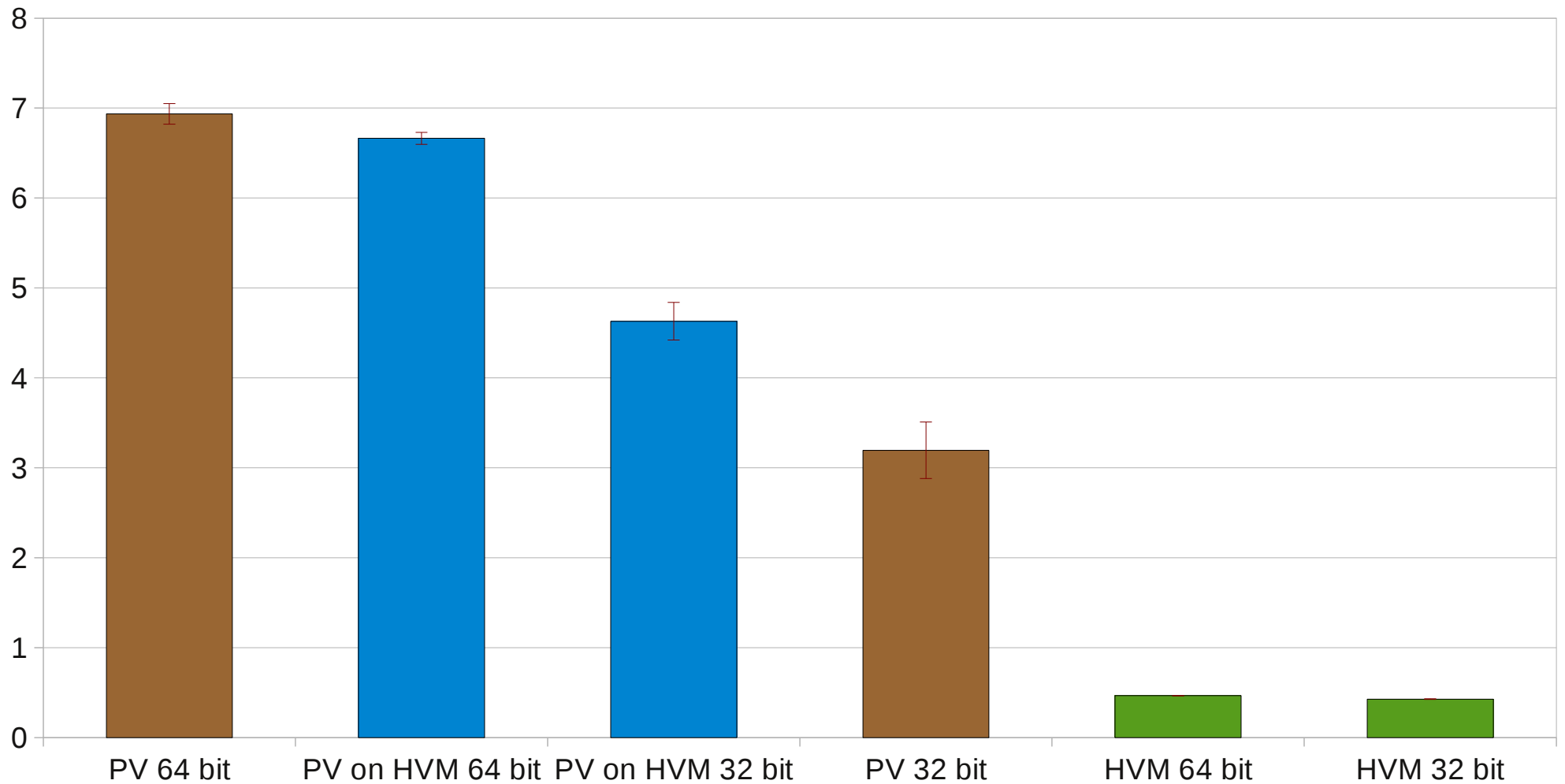
# SPECjbb2005

Results: percentage of native, the higher the better



# Iperf tcp

Results: gbit/sec, the higher the better



# Conclusions

PV on HVM guests are very close to PV guests in benchmarks that favor PV MMUs

PV on HVM guests are far ahead of PV guests in benchmarks that favor nested paging

Questions?